**Directions:** Solve 5 of the following 6 problems. All written work must be your own, using only permitted sources. See the "General Guidelines and Advice" on the homework page for more details.

- 1. [IGT 1.3.22] Let G be a nonbipartite triangle-free graph with n vertices and minimum degree k. Let l be the minimum length of an odd cycle in G.
  - (a) Let C be a cycle of length l in G. Prove that every vertex not in V(C) contains at most two neighbors in V(C).
  - (b) By counting the edges joining V(C) and V(G) V(C) in two ways, prove that  $n \ge kl/2$  (and thus  $l \le 2n/k$ ).
  - (c) When k is even, prove that the inequality of part (b) is best possible. (Hint: form a graph having k/2 pairwise disjoint *l*-cycles.)
- 2. [IGT 1.3.44] Let G be a graph with average degree a. Recall that a = 2|E(G)|/|V(G)|.
  - (a) Prove that G x has average degree at least a if and only if  $d(x) \le a/2$ .
  - (b) Prove that if a > 0, then G has a subgraph with minimum degree greater than a/2. Hint: what happens if we delete a vertex of degree at most a/2?
  - (c) Show that there is no constant c greater than 1/2 such that G must have a subgraph with minimum degree greater than ca; this proves that the bound in part (b) is best possible. Hint: Use  $K_{1,n-1}$ .
- 3. [IGT 1.4.25]
  - (a) Prove that every connected graph has an orientation in which the number of vertices with odd outdegree is at most 1. (Hint: consider an orientation with the fewest number of vertices with odd outdegree.)
  - (b) Use part (a) to conclude that every connected graph with an even number of edges has a  $P_3$ -decomposition.
- 4. [IGT 1.4.29] Suppose that G is a graph and D is an orientation of G that is strongly connected. Prove that if G has an odd cycle, then D has a (directed) odd cycle. (Hint: consider each pair  $\{v_i, v_{i+1}\}$  in an odd cycle  $v_1 \cdots v_k$  of G.)
- 5. [IGT 1.4.34] Let G and H be tournaments on the same vertex set. Prove that  $d_G^+(v) = d_H^+(v)$  for each vertex v if and only if G can be turned into H by a sequence of direction-reversals on cycles of length 3.
- 6. Let G be a directed graph without loops. Prove that G has an independent set S such that every vertex in G is reachable from a vertex in S by a directed path of length at most 2. Hint: use induction on |V(G)| and recall that the induction hypothesis applies to all graphs with fewer vertices, not just graphs with |V(G)| - 1 vertices.