

Ordered Turán Numbers

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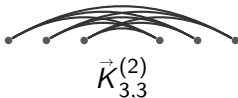
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- ▶ Note: $\vec{K}_{1,2}^{(2)}, \vec{K}_{2,1}^{(2)} \subseteq \vec{K}_{3,3}^{(2)}$ but $\vec{P}_3^{(2)} \not\subseteq \vec{K}_{3,3}^{(2)}$.

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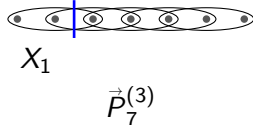


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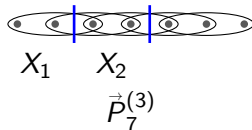
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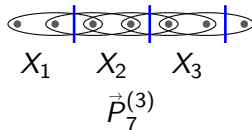
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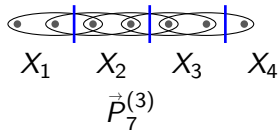
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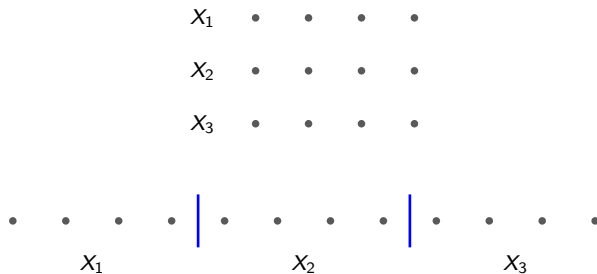
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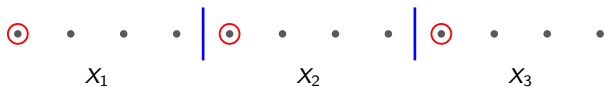
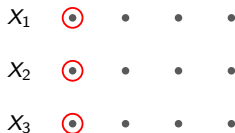
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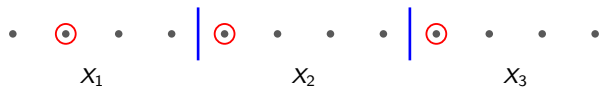
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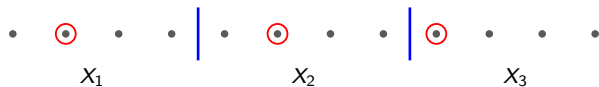
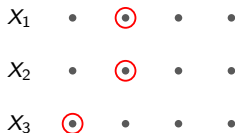
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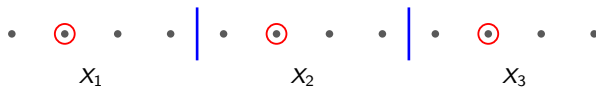
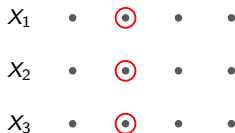
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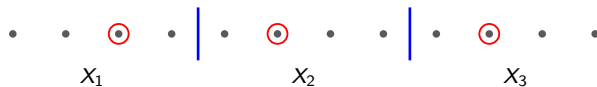
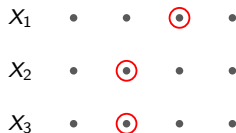
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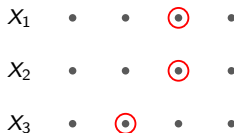
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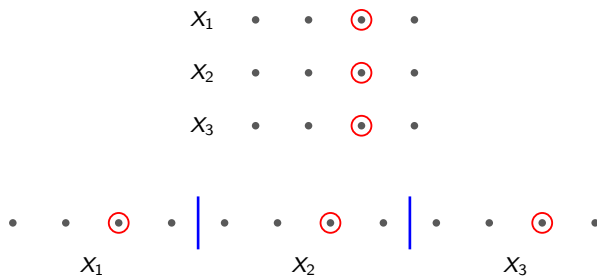
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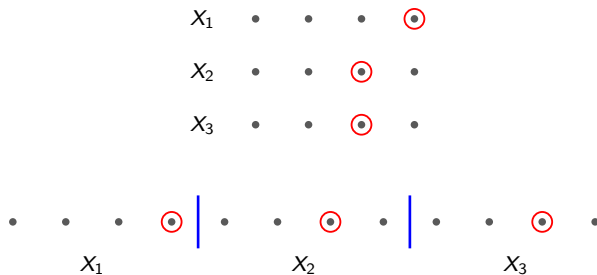
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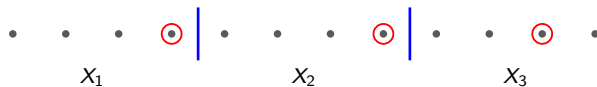
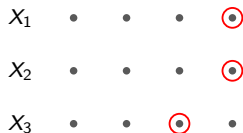
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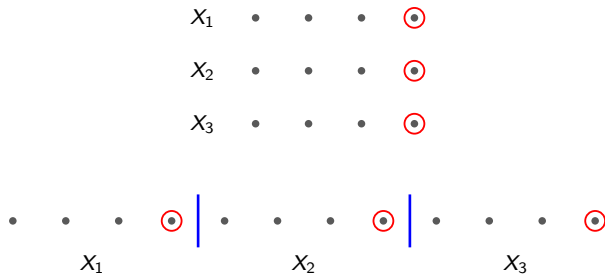
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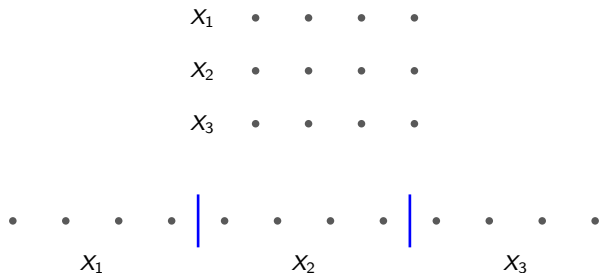
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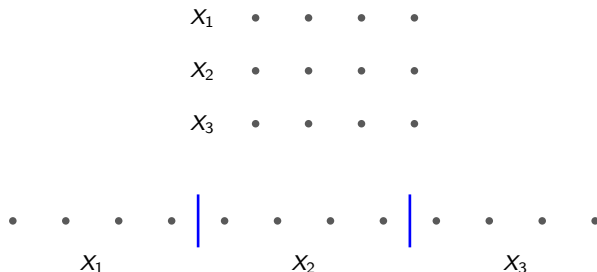


- ▶ Thm (Füredi–Jiang–Kostochka–Mubayi–Verstraëte 2021):

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- ▶ Complementary transversal numbers for r -uniform H :

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Let n be even, let $r \leq s \leq 2r - 1$, let $m = s - r + 1$, and let $\vec{LP}_s^{(r)}$ be the *loose path* obtained from $\vec{P}_s^{(r)}$ by removing all except the first and last edges. Let $\alpha = 2h(n, r, m) + h(n, r - 1, m)$. We have

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- ▶ Define **m -right-biased** in the natural way. A set is **m -biased** if it is m -left-biased or m -right-biased.

Loose path transversals



Theorem

Let $r \leq s \leq 2r - 1$, let $m = s - r + 1$, and let n be even. We have $\vec{\tau}(n, \vec{LP}_s^{(r)}) \leq 2h(n, r, m) + h(n, r - 1, m)$.

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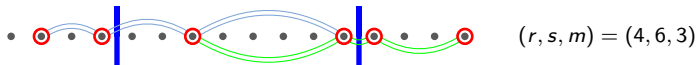
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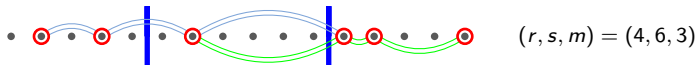
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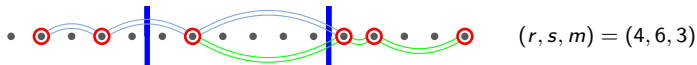
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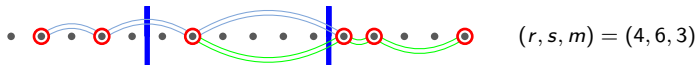
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- ▶ Such a partition exists as otherwise $s \leq 2(m - 1) = 2(s - r)$ and so $s \geq 2r$.



Loose path transversals



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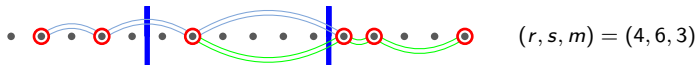
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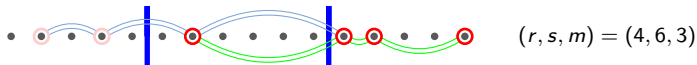
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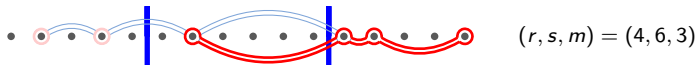
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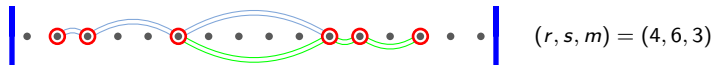
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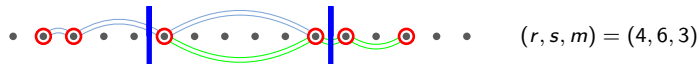
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Tight path packings

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Let $r \leq s \leq 2r - 1$, let $m = s - r + 1$, and let n be even. We have

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- ▶ Start with the transversal graph $G \subseteq \vec{K}_n^{(r)}$.
- ▶ For each $e \in E(G)$, construct a copy Q_e of $\vec{P}_s^{(r)}$.
- ▶ Show that $\{Q_e : e \in E(G)\}$ is edge-disjoint.

Fractional Variants

Frac. H packing $\vec{v}^*(n, H)$:

$$\begin{aligned} \text{max:} & \sum_{H \subseteq \vec{K}_n^{(r)}} y_H \\ \text{subj to: } \forall e & \sum_{H: e \in E(H)} y_H \leq 1 \end{aligned}$$

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 - ▶ We need a new approach for large s , probably even $s = 2r$.

The case $r = 3$

► Thm: $\vec{\tau}(n, \vec{P}_s^{(3)}) \leq \left(\frac{1}{\lfloor \frac{(s-1)^2}{4} \rfloor} + o(1) \right) \binom{n}{3}$

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- ▶ Let $f_k(n)$ be the max., over all $[k]$ -edge-labelings of $\vec{K}_n^{(2)}$, of the number of triples $u < v < w$ with $\ell(uv) < \ell(vw)$.

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- ▶ Prop: $\vec{e}x(n, \vec{P}_s^{(3)}) = f_{s-2}(n)$

Increasing Triple Optimization

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Increasing Triple Optimization

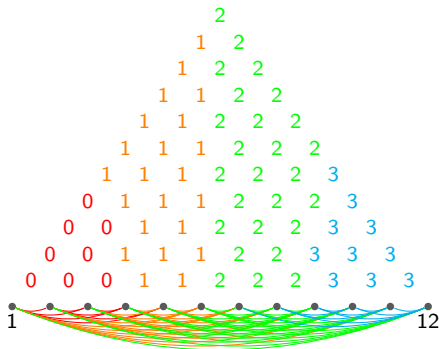
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- ▶ First open case is $\vec{P}_6^{(3)}$, with 4 edge labels.

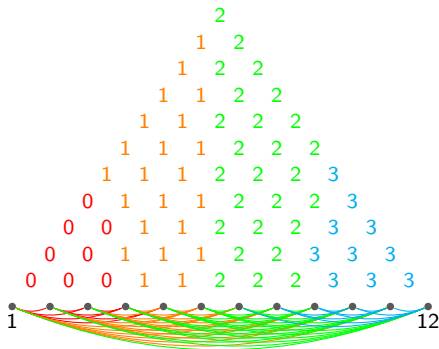
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- ▶ Prop: $\vec{ex}(n, \vec{P}_5^{(3)}) = f_{5-2}(n)$.
- ▶ First open case is $\vec{P}_6^{(3)}$, with 4 edge labels.
- ▶ Conjecture: $f_4(n) = \vec{ex}(n, \vec{P}_6^{(3)}) = \left(\frac{5}{6} + o(1)\right) \binom{n}{3}$



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Thank You.