## First-Fit chain partitions in partially ordered sets

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- In terms of the width $w$, how many chains can First-Fit use?

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Theorem (Bosek-Krawczyk-Matecki (2011))
If $Q$ has width at most 2 , then $\operatorname{FF}(w, Q)$ is bounded.

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Main Question
What properties of $Q$ determine the behavior of $\operatorname{FF}(w, Q)$ ?

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- Similarly $\max W_{i}>\min W_{j}$.
- Let $y \in W_{j}$, and let $x \in W_{i}$ be incomparable to $y$.
- Note $\min W_{i}<x, y<\max W_{i}$, completing a copy of $L_{2}$.


## Ladders



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- Bosek-Krawczyk (2015): if $\mathrm{FF}\left(w, L_{m}\right)$ is polynomial in $w$ and $m$, then there is a polynomial online chain partitioning algorithm.


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- The upper bound is superpolynomial but subexponential.


## Our results

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Theorem (First-Fit Dichotomy)

- If $Q \in \mathcal{Q}$, then $\mathrm{FF}(w, Q) \leq w^{c_{Q} \log w}$ for some constant $c_{Q}$.


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1. If $Q$ is ladder-like, then $Q \in \mathcal{Q}$.
2. If $Q_{1}, Q_{2} \in \mathcal{Q}$, then $Q_{1} \otimes Q_{2} \in \mathcal{Q}$.

Theorem (First-Fit Dichotomy)

- If $Q \in \mathcal{Q}$, then $\mathrm{FF}(w, Q) \leq w^{c_{Q} \log w}$ for some constant $c_{Q}$.
- If $Q \notin \mathcal{Q}$, then $\operatorname{FF}(w, Q) \geq 2^{w}-1$.


## First-Fit Dichotomy Theorem, Upper Bound

Proposition
If $Q$ is a ladder-like m-point poset, then $Q$ is a subposet of $L_{m}$.

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- Therefore $\operatorname{FF}(w, Q) \leq w^{c} Q^{\log w}$ when $Q \in \mathcal{Q}$.


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- First-Fit uses $2 s+1$ colors, and $2 s+1=2^{w}-1$.
- Prop: a poset $Q$ of width 2 is in $\mathcal{Q}$ if and only if $Q$ is a subposet of some $R_{w}$.
- So, if $Q \notin \mathcal{Q}$, then $\operatorname{FF}(w, Q) \geq 2^{w}-1$.


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Skewed butterfly $\widehat{B}$

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- On the other hand, $\widehat{B} \notin \mathcal{Q}$ and so $\mathrm{FF}(w, \widehat{B}) \geq 2^{w}-1$.


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Thank You.

