

## Covering the Hypercube with a Bounded Number of Disjoint Snakes

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**Abstract.** We present a construction of an induced cycle in the  $n$ -dimensional hypercube  $I[n]$  ( $n \geq 2$ ), and a subgroup  $\mathcal{H}_n$  of  $I[n]$  considered as the group  $\mathbb{Z}_2^n$ , such that  $|\mathcal{H}_n| \leq 16$  and the induced cycle uses exactly one element of every coset of  $\mathcal{H}_n$ . This proves that for any  $n \geq 2$  the vertices of  $I[n]$  can be covered using at most 16 vertex-disjoint induced cycles.

## 1. Introduction

Given a positive integer  $d$ , let  $I[d]$  be the  $d$ -dimensional cube (also called hypercube) *i.e.* the graph with all  $d$ -tuples of binary digits as vertices, and all pairs of vertices differing at exactly one coordinate as edges. It will be convenient to think of the set of vertices of  $I[d]$  as the set of elements of the group  $\mathbb{Z}_2^d = (\mathbb{Z}/2\mathbb{Z})^d$  with the operation  $\oplus$  of componentwise addition mod 2. Thus, when we refer to  $I[d]$ , we assume that it has both a structure of a graph and a structure of a group.

A *snake* is an induced cycle in  $I[d]$ . For each  $d \geq 2$ , let  $S(d)$  denote the length of the longest snake in  $I[d]$ . An extensive literature has evolved concerning the problem of estimating  $S(d)$ . See [1], [3], [5], and the references in these papers. What we know now is that

$$\frac{77}{256}2^d \leq S(d) \leq 2^{d-1} - \frac{2^{d-1}}{20d - 41},$$

assuming that  $d \geq 12$  in the upper bound. The lower bound was proved by Abbott and Katchalski [2] and the upper bound by Snevily [4].

During the XXIII Southeastern International Conference, Boca Raton 1992, Erdős posed the problem of deciding whether there is a number  $k$  such that for every  $d \geq 2$  the vertices of  $I[d]$  can be covered using at most  $k$  snakes, and if the answer to the above problem is positive, then whether it can be done in such a way that the snakes are pairwise vertex-disjoint. In this note we show that the answer to both of the above questions is positive with  $k = 16$ . Actually, we prove the following theorem, which is a stronger result.

**Theorem 1.** *For every  $n \geq 2$ , there is a subgroup  $\mathcal{H}_n \subset I[n]$  and a snake  $C_n \subset I[n]$  such that  $|\mathcal{H}_n| \leq 16$  and  $C_n$  uses exactly one element of every coset of  $\mathcal{H}_n$ .*

## 2. Basic Definitions

Let  $\psi_0, \psi_1 : I[d] \rightarrow I[d+1]$  be the embeddings defined by

$$\psi_i(v_1 v_2 \dots v_d) = v_1 v_2 \dots v_d i,$$

for  $i = 0, 1$ . If  $F$  is a subgraph of  $I[d]$ , let  $F^{(i)}$  be the subgraph of  $I[d+1]$  obtained as the image of  $F$  under the embedding  $\psi_i$ .

For each  $d \geq 2$  we define a function  $H_d : \{1, 2, \dots, 2^d\} \rightarrow V(I[d])$  such that  $H_d^* = (H_d(1), \dots, H_d(2^d), H_d(1))$  is a Hamiltonian cycle in  $I[d]$ . We set

$$H_2^* = (00, 01, 11, 10, 00),$$

and

$$H_{d+1}(i) = \begin{cases} (H_d(i))^{(0)} & \text{if } 1 \leq i \leq 2^d, \\ (H_d \circ R_d(i))^{(1)} & \text{if } 2^d + 1 \leq i \leq 2^{d+1}, \end{cases}$$

where  $R_d : \{2^d + 1, 2^d + 2, \dots, 2^{d+1}\} \rightarrow \{1, 2, \dots, 2^d\}$  is the order reversing bijection,

$$R_d(i) = 2^{d+1} + 1 - i.$$

In other words,  $H_{d+1}^*$  is obtained by taking  $H_d^{*(0)}$  and  $H_d^{*(1)}$ , removing the edges connecting their last vertices with their first vertices, joining the first vertex of  $H_d^{*(0)}$  with the first vertex of  $H_d^{*(1)}$  and analogously the last with the last.

Let us regard  $I[d+6]$  as  $I[d] \times I[6]$ , that is as the  $d$ -dimensional cube  $I[d]$  with each vertex being a copy of  $I[6]$ . Suppose that  $P_d = \{(v_j^1, \dots, v_j^{r_j})\}_{j=1}^{2^d}$  is a sequence of  $2^d$  paths in  $I[6]$  such that  $v_i^{r_i} = v_j^1$  when  $1 \leq i \leq 2^d - 1$  and  $j = i + 1$  or  $i = 2^d$  and  $j = 1$ . Such a sequence will be called a  $2^d$ -chain of paths in  $I[6]$ . We can use  $P_d$  and  $H_d^*$  to construct a cycle  $C_{d+6}$  in  $I[d+6]$ . Let us take the  $j$ th path  $(v_j^1, \dots, v_j^{r_j})$  in the copy of  $I[6]$  corresponding to the vertex  $H_d(j)$  in  $I[d]$ ; see Figure 1 for the case  $d = 3$ .

Then, let us join  $v_i^{r_i}$  from the  $i$ th copy of  $I[6]$  with  $v_{i+1}^1$  from the  $(i+1)$ st copy of  $I[6]$  for all  $i \in \{1, \dots, 2^d\}$ , where the indices are understood circularly. Hence we have

$$C_{d+6} = ((H_d(1), v_1^1), (H_d(1), v_1^2), \dots, (H_d(1), v_1^{r_1}), (H_d(2), v_2^1), \dots, (H_d(2), v_2^{r_2}), \dots, \\ (H_d(2^d), v_{2^d}^1), \dots, (H_d(2^d), v_{2^d}^{r_{2^d}}), (H_d(1), v_1^1)).$$

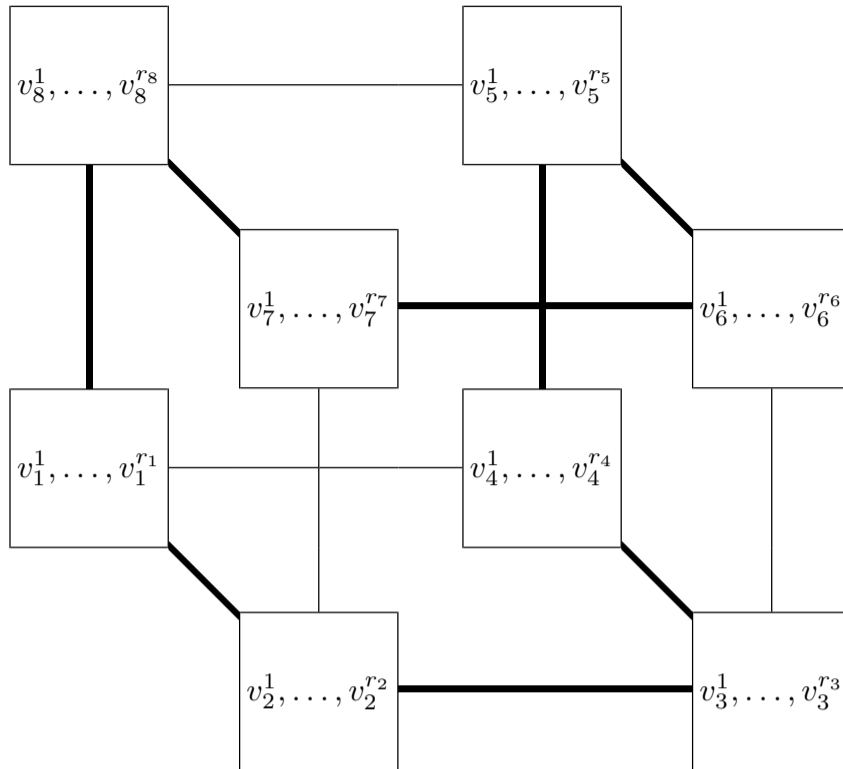


Fig. 1.

It is clear that  $C_{d+6}$  is a cycle in  $I[d+6]$ , we will call it the cycle *generated* by  $H_d$  and  $P_d$ . We call  $P_d$  *well separated* with respect to  $H_d$  if the cycle  $C_{d+6}$  generated by  $H_d$  and  $P_d$  is an induced cycle.

Let  $\mathcal{P}$  be the set of all 12 paths in  $I[6]$  that can be obtained from one of the following two paths

$$p_0 = (000000, 100000, 110000, 111000)$$

$$q_0 = (111000, 111100, 111110, 111111)$$

by a cyclic permutation of coordinates of every vertex of the path. Let  $\mathcal{H}$  be the subgroup of  $I[6]$  containing all  $a_1 a_2 \dots a_6 \in I[6]$  such that the following two conditions are satisfied

- (i)  $a_1 + a_3 + a_5$  is even, and
- (ii)  $a_2 + a_4 + a_6$  is even.

Note that  $\mathcal{H}$  is generated by the set of elements of  $I[6]$  that can be obtained from

101000 by cyclic permutations of the coordinates. Also note that the elements of  $\mathcal{H}$  are 000000, six cyclic permutations of 101000, six cyclic permutations of 111100, and three cyclic permutations of 110110, so  $|\mathcal{H}| = 16$ . To prove Theorem 1, we will define a  $2^d$ -chain  $P_d$  of paths which is well separated with respect to  $H_d$  and the paths in  $P_d$  are elements of  $\mathcal{P}$ . We will need the following lemma.

**Lemma 1.** *Every path  $P \in \mathcal{P}$  uses exactly one element of each coset of  $\mathcal{H}$  in  $I[6]$ .*

*Proof.* Let  $P \in \mathcal{P}$ , and let  $v_1, v_2$  be two vertices of  $P$ . It is easy to observe that if  $v_1 \neq v_2$  then  $v_1 \oplus v_2$  has an odd number of ones or two consecutive ones (in the cyclic order), where  $\oplus$  is the operation of the group  $I[6]$ . Thus, if  $v_1 \oplus v_2 \in \mathcal{H}$  then  $v_1 = v_2$ .

This proves that  $P$  uses at most one element of every coset of  $\mathcal{H}$ . Since  $|P| \times |\mathcal{H}| = |I[6]|$ , the path  $P$  uses exactly one element of every coset of  $\mathcal{H}$ , and the proof is complete. □

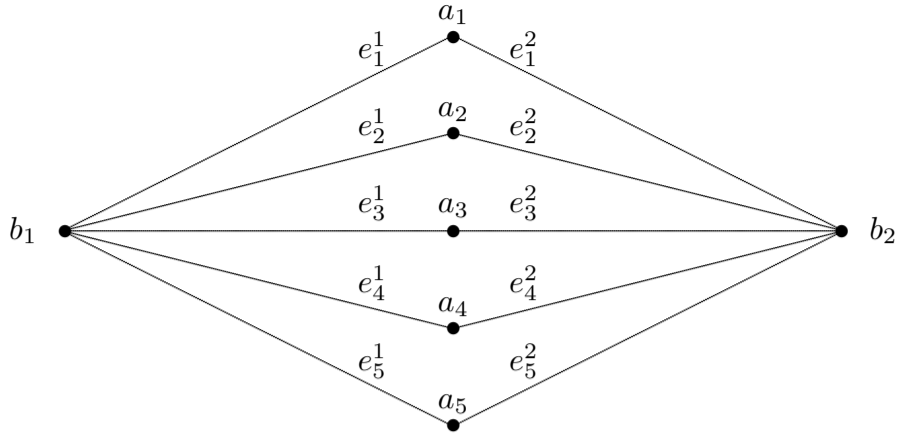


Fig. 2. The graph  $G$

### 3. The Main Result

Let  $G = K_{2,5}$  be the graph shown in Figure 2.

The following lemma is proved in Wojciechowski [5] (Lemma 3).

**Lemma 2.** For every  $d \geq 2$  there is a function  $\Phi_d : \{1, 2, \dots, 2^d\} \rightarrow E(G)$  such that

(i) if  $1 \leq i \leq 2^d - 1$  and  $j = i + 1$ , or else  $i = 2^d$ ,  $j = 1$ , then  $\Phi_d(i)$  and  $\Phi_d(j)$  have exactly one vertex  $v_i$  in common, such that  $v_i \in \{a_1, \dots, a_5\}$  for  $i$  even, and  $v_i \in \{b_1, b_2\}$  for  $i$  odd, and

(ii) if  $(H_d(i), H_d(j)) \in E(I[d]) \setminus E(H_d^*)$ , then  $\Phi_d(i)$  and  $\Phi_d(j)$  are vertex-disjoint.  $\square$

Let  $\mathcal{P}$  be the set of paths defined in Section 2. The following lemma will be used in the proof of the main result.

**Lemma 3.** For every  $d \geq 2$  there is a  $2^d$ -chain  $P_d$  of paths in  $I[6]$  such that every path in  $P_d$  belongs to  $\mathcal{P}$ , and  $P_d$  is well separated with respect to  $H_d$ .

*Proof.* Let  $G'$  be the subdivision of  $G$  obtained by subdividing  $e_k^1$  ( $1 \leq k \leq 5$ ) with two new vertices  $c_k^1$  and  $c_k^2$  in such a way that we get the path  $(b_1, c_k^1, c_k^2, a_k)$ , and subdividing  $e_k^2$  with two new vertices  $c_k^4, c_k^5$ , giving rise to the path  $(a_k, c_k^4, c_k^5, b_2)$ . Let  $c_k^3 = a_k$  and

$\xi : V[G'] \rightarrow V(I[6])$  be defined as follows. Set

$$\xi(b_1) = 000000,$$

$$\xi(c_1^1) = 100000,$$

$$\xi(c_1^2) = 110000,$$

$$\xi(c_1^3) = 111000,$$

$$\xi(c_1^4) = 111100,$$

$$\xi(c_1^5) = 111110,$$

$$\xi(b_2) = 111111,$$

and if  $\xi(c_1^i) = \alpha_1 \dots \alpha_6$ , then let

$$\xi(c_k^i) = \alpha_k \dots \alpha_6 \alpha_1 \dots \alpha_{k-1},$$

for  $2 \leq k \leq 5$ .

It is clear that the function  $\xi$  defines an embedding of  $G'$  into  $I[6]$  such that the image of the subdivision of any edge of  $G$  is an induced path in  $I[6]$ . Let  $\Phi_d : \{1, 2, \dots, 2^d\} \rightarrow E(G)$  be a function satisfying conditions (i) and (ii) of Lemma 2. Let  $P_d = \{(v_i^1, \dots, v_i^4)\}_{i=1}^{2^d}$  be a  $2^d$ -chain of paths in  $I[6]$  such that  $(v_i^1, \dots, v_i^4)$  is the image under  $\xi$  of the subdivision of the edge  $\Phi_d(i)$ .

Let  $C_{d+6}$  be the cycle generated by  $H_d$  and  $P_d$ . Since each pair of vertex-disjoint edges in  $G$  corresponds to a pair of vertex-disjoint paths in  $P_d$ , and since each pair of edges having exactly one vertex in common corresponds to a pair of paths in  $P_d$  having exactly one vertex in common, it follows from (i) and (ii) of Lemma 2 that  $C_{d+6}$  is an induced cycle. So  $P_d$  is well separated with respect to  $H_d$ , and the proof is complete.  $\square$

We can now prove our main result.

*Proof of Theorem 1.* Let us assume first that  $n \geq 8$ . Set  $d = n - 6$ . Let  $P_d$  be a  $2^d$ -chain of paths in  $I[6]$  such that every path in  $P_d$  belongs to  $\mathcal{P}$  and  $P_d$  is well separated with

respect to  $H_d$ , and let  $C_n$  be the snake generated by  $P_d$  and  $H_d$ . Let  $\mathcal{H}_n$  be the subgroup of  $I[n]$  defined to be the set of all  $(a_1, a_2, \dots, a_d, b_1, \dots, b_6) \in I[n]$  such that the following conditions are satisfied.

- (i)  $a_1 = a_2 = \dots = a_d = 0$ ,
- (ii)  $b_1 + b_3 + b_5$  is even, and
- (iii)  $b_2 + b_4 + b_6$  is even.

We will show that  $C_n$  contains exactly one element of every coset of  $\mathcal{H}_n$ . Since  $|C_n| = 4 \times 2^d$  and  $|\mathcal{H}_n| = 16$ , it is enough to prove that  $C_n$  contains at most one element of every coset of  $\mathcal{H}_n$ . Suppose  $w_1, w_2 \in C_n$  and the cosets of  $w_1$  and  $w_2$  are equal. Then  $w_1 \oplus w_2 \in \mathcal{H}_n$ . Let  $v_1 \in I[6]$  be the sequence of the last six digits of  $w_1$ , and similarly let  $v_2 \in I[6]$  consist of the last six digits of  $w_2$ . Since  $w_1 \oplus w_2 \in \mathcal{H}_n$ , the first  $d$  digits of  $w_1$  are the same as the first  $d$  digits of  $w_2$ , and  $v_1 \oplus v_2 \in \mathcal{H}$ . Thus  $v_1$  and  $v_2$  are vertices of the same path in  $\mathcal{P}$ , and so it follows from Lemma 1 that  $v_1 = v_2$ . Hence  $w_1 = w_2$ , and the proof of the case  $n \geq 8$  is complete.

If  $2 \leq n \leq 6$ , then we can take  $C_n = (000 \dots 0, 100 \dots 0, 110 \dots 0, 010 \dots 0, 000 \dots 0)$ , and  $\mathcal{H}_n$  to be the set of all elements of  $I[n]$  with first two coordinates equal 0. It is clear that  $C_n$  uses exactly one element of every coset of  $\mathcal{H}_n$ . In the remaining case  $n = 7$ , let  $C_7 = (0000000, 1000000, 1100000, 1110000, 1111000, 0111000, 0011000, 0001000, 0000000)$ , and let  $\mathcal{H}_7$  be the set of all elements of  $I[7]$  with the first four coordinates being either 0000 or 1010. It is straightforward to check that in this case the conclusion is true as well. Thus the proof is complete.  $\square$

The following corollary gives the answer to the problem of Erdős.

**Corollary 1.** *For every  $n \geq 2$  the vertices of  $I[n]$  can be covered using at most 16 pairwise disjoint snakes.*

*Proof.* Let us take a subgroup  $\mathcal{H}_n$  and a snake  $C_n$  in  $I[n]$  such that  $|\mathcal{H}_n| \leq 16$  and  $C_n$



uses exactly one element of every coset of  $\mathcal{H}_n$ . The family of snakes  $\{C_n \oplus h : h \in \mathcal{H}_n\}$  contains 16 vertex-disjoint snakes covering  $I[n]$ .  $\square$

#### 4. Concluding remarks

Let  $k_0$  be the smallest integer such that for every  $n \geq 2$ , the cube  $I[n]$  can be vertex covered by at most  $k_0$  snakes. Let  $k_1$  and  $k_2$  be defined in a similar way taking pairwise vertex-disjoint snakes and pairwise vertex-disjoint snakes of equal length, respectively. Set  $k_3$  to be the smallest integer such that for every  $n \geq 2$  there is a subgroup  $\mathcal{H}_n$  and a snake  $C_n$  in  $I[n]$  such that  $|\mathcal{H}_n| \leq k_3$  and  $C_n$  uses exactly one element of every coset of  $\mathcal{H}_n$ . As a corollary of Theorem 1 and the upper bound for the length of snakes we get the following theorem.

**Theorem 2.** *We have  $3 \leq k_0 \leq k_1 \leq k_2 \leq k_3 \leq 16$  and  $k_2, k_3 \in \{4, 8, 16\}$ .*  $\square$

The question of determining the exact values of  $k_0$ ,  $k_1$ ,  $k_2$ , and  $k_3$  remains open. It might be possible to modify the technique used in this note to improve Theorem 2 to make the upper bound equal 8 or perhaps even 4. The possible approach may involve finding a more sophisticated embedding of a subdivision of  $K_{2,5}$  into  $I[c]$ , where  $c$  is a small integer constant (in our proof we used  $c = 6$ ), or else replacing  $K_{2,5}$  by another graph, perhaps  $K_{2,4}$ .

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