

Note

# Quadrangulary connected claw-free graphs<sup>☆</sup>

MingChu Li<sup>a</sup>, Cheng Guo<sup>a</sup>, Liming Xiong<sup>b</sup>, Dengxin Li<sup>c</sup>, Hong-Jian Lai<sup>d</sup>

<sup>a</sup>School of Software, Dalian University of Technology, Dalian 116024, PR China

<sup>b</sup>Department of Mathematics, Beijing Institute of Technology, Beijing 100081, PR China

<sup>c</sup>Department of Mathematics, Chongqing Technology and Business University, Chongqing 400067, PR China

<sup>d</sup>Department of Mathematics, West Virginia University, Morgantown, WV 26506, USA

Received 27 November 2004; received in revised form 12 June 2006; accepted 20 July 2006

Available online 10 October 2006

## Abstract

A graph  $G$  is *quadrangulary connected* if for every pair of edges  $e_1$  and  $e_2$  in  $E(G)$ ,  $G$  has a sequence of  $l$ -cycles ( $3 \leq l \leq 4$ )  $C_1, C_2, \dots, C_r$  such that  $e_1 \in E(C_1)$  and  $e_2 \in E(C_r)$  and  $E(C_i) \cap E(C_{i+1}) \neq \emptyset$  for  $i = 1, 2, \dots, r - 1$ . In this paper, we show that every quadrangulary connected claw-free graph without vertices of degree 1, which does not contain an induced subgraph  $H$  isomorphic to either  $G_1$  or  $G_2$  such that  $N_1(x, G)$  of every vertex  $x$  of degree 4 in  $H$  is disconnected is hamiltonian, which implies a result by Z. Ryjáček [Hamiltonian circuits in  $N_2$ -locally connected  $K_{1,3}$ -free graphs, J. Graph Theory 14 (1990) 321–331] and other known results.

© 2006 Elsevier B.V. All rights reserved.

**Keywords:** Cycle; Claw-free graph; Quadrangulary connected

## 1. Notation and terminology

We use [1] for notation and terminology not defined here, and consider finite simple graphs only. Let  $G$  be a graph. Denote by  $G[S]$  the induced subgraph of  $G$  on the subset  $S$  of  $V(G)$ . For a vertex  $v$  of  $G$ , the neighborhood of  $v$ , i.e., the induced subgraph on the set of all vertices that are adjacent to  $v$ , will be called the neighborhood of the *first type* of  $v$  in  $G$  and denoted by  $N_1(v, G)$ , or briefly,  $N(v)$ . For notational convenience, we shall use  $N_G(v)$  to denote both the induced subgraph and the set of vertices adjacent to  $v$  in  $G$ . We define the neighborhood of the *second type* of  $v$  in  $G$  (denoted by  $N_2(v, G)$ , or briefly,  $N_2(v)$ ) as the subgraph of  $G$  induced by the edge subset  $\{e = xy \in E(G) : v \notin \{x, y\} \text{ and } \{x, y\} \cap N(v) \neq \emptyset\}$ . We say that a vertex  $v$  is *locally connected* if  $N(v)$  is connected; and  $G$  is *locally connected* if every vertex of  $G$  is locally connected. Analogously, a vertex  $v$  is  *$N_2$ -locally connected* if its second type neighborhood  $N_2(v)$  is connected; and  $G$  is called  *$N_2$ -locally connected* if every vertex of  $G$  is  $N_2$ -locally connected. It follows from the definitions that every locally connected graph is  $N_2$ -locally connected. A cycle of length  $k$  is called a  $k$ -cycle. Given a cycle  $C = (x_1x_2 \dots x_kx_1)$  of a graph  $G$ , we fix an orientation of  $C$ , let  $x_i^+ = x_{i+1}$  and  $x_i^- = x_{i-1}$ , and let  $C[x_i, x_j] = (x_ix_{i+1}, \dots, x_{j-1}, x_j) = \{x_i, x_{i+1}, \dots, x_{j-1}, x_j\}$ ,  $C(x_i, x_j) = C[x_i, x_j] - \{x_i, x_j\}$  and

<sup>☆</sup> Supported by Nature Science Foundation of China under grants no.: 90412007 (M. Li), 60673046 (M. Li), and 10671014 (L. Xiong) and by Excellent Young Scholar Research Fund of Beijing Institute of Technology (No.: 000Y07-28) (L. Xiong).

E-mail address: [li\\_minghu@yahoo.com](mailto:li_minghu@yahoo.com) (M. Li).

$C^-[x_j, x_i] = (x_j x_{j-1} \dots x_i)$ . An edge  $e = uv$  is called a chord of  $C$  if  $e \notin E(C)$  and  $u$  and  $v$  are on  $C$ . If a cycle  $C$  has no chords, then we say that  $C$  is chord-free. Obviously, 3-cycles are chord-free.

The line graph of a graph  $G$ , denote it by  $L(G)$ , has  $E(G)$  as its vertex set, where two vertices in  $L(G)$  are adjacent if and only if the corresponding edges in  $G$  are adjacent. A graph  $G$  is *claw-free* if it does not contain  $K_{1,3}$  as an induced subgraph. Obviously, the line graph of a graph is claw-free.

**2. Introduction**

The following theorem gives the hamiltonicity of locally connected graphs.

**Theorem 2.1** (Oberly and Sumner [6]). *Every connected locally connected claw-free graph on at least three vertices is hamiltonian.*

A pair of edges  $e_1$  and  $e_2$  in  $E(G)$  is called *triangularly adjacent* if  $G$  has a sequence of 3-cycles  $C_1, C_2, \dots, C_r$  such that  $e_1 \in E(C_1)$  and  $e_2 \in E(C_r)$  and  $E(C_i) \cap E(C_{i+1}) \neq \emptyset$  for  $i = 1, 2, \dots, r - 1$ . A graph  $G$  is *triangularly connected* if every pair of edges  $e_1$  and  $e_2$  in  $E(G)$  is triangularly adjacent. Obviously, every connected, locally connected graph is triangularly connected (see [2]). But the converse is not true. Recently, Shao [2] generalized the above theorem as follows.

**Theorem 2.2** (Shao [2]). *Every triangularly connected claw-free graph on at least three vertices is hamiltonian.*

A graph  $G$  is *vertex pancyclic* if it contains cycles of all possible length through every vertex. They [2] actually proved that a graph with the same conditions as Theorem 2.2 is vertex pancyclic. Ryjáček [7] strengthened Theorem 2.1 as follows in 1990, and Li [5] improved Theorem 2.3 using  $N^2$ -locally connectedness.

**Theorem 2.3** (Ryjáček [7]). *Let  $G$  be a connected,  $N_2$ -locally connected claw-free graph without vertices of degree 1, which does not contain an induced subgraph  $H$  isomorphic to either  $G_1$  or  $G_2$  (Fig. 1) such that  $N_1(x, G)$  of every vertex  $x$  of degree 4 in  $H$  is disconnected. Then  $G$  is hamiltonian.*

Li [4] obtained the following theorem.

**Theorem 2.4** (Li [4]). *Let  $G$  be a connected,  $N_2$ -locally connected claw-free graph with  $\delta(G) \geq 3$ , which does not contain an induced subgraph  $H$  isomorphic to either  $G_1$  or  $G_2$  (Fig. 1). Then  $G$  is vertex pancyclic.*

A pair of edges  $e_1$  and  $e_2$  in  $E(G)$  is called *quadrangularly connected* if  $G$  has a sequence of chord-free  $l$ -cycles ( $3 \leq l \leq 4$ )  $C_1, C_2, \dots, C_r$  such that  $e_1 \in E(C_1)$  and  $e_2 \in E(C_r)$  and  $E(C_i) \cap E(C_{i+1}) \neq \emptyset$  for  $i = 1, 2, \dots, r - 1$ .

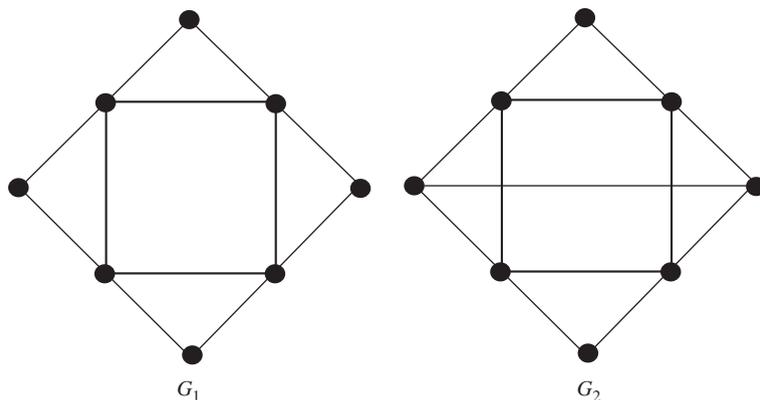


Fig. 1.

A graph  $G$  is *quadrangularly connected* if every pair of edges in  $E(G)$  is quadrangularly adjacent. Obviously, if a graph is *triangularly connected* then it is quadrangularly connected. From the definition, we easily prove the following proposition that every connected  $N_2$ -locally connected graph is quadrangularly connected. But the converse is not true. For example, let  $k > 4$  be an integer, and let  $H_1, H_2, \dots, H_k$  be complete graphs of order at least 4. We construct a graph  $G$  by the following process: for each  $i = 1, 2, \dots, k - 1$ , choose an edge  $e'_i = v'_{i1}v'_{i2}$  in  $H_i$  and an edge  $e''_i = v''_{i1}v''_{i2}$  in  $H_{i+1}$ , and add two new edges  $e_i^1 = v'_{i1}v''_{i1}$  and  $e_i^2 = v'_{i2}v''_{i2}$  to form a quadrangle of  $G$ , and identify a vertex  $v_1$  in  $H_1$  with a vertex  $v_k$  in  $H_k$  to form a vertex  $v$  in  $G$  such that  $V(H_1) \cap V(H_k) = \{v\}$  and  $v \notin \{v'_{11}, v'_{12}, v''_{(k-1)1}, v''_{(k-1)2}\}$ . Then  $G$  is quadrangularly connected, but  $v$  is not  $N_2$ -locally connected because  $N_2(v)$  does not induce a connected subgraph of  $G$ . Note that  $G$  is not triangularly connected because  $e'_1$  and  $e''_1$  are not triangularly adjacent.

Define a relation on  $E(G)$  such that  $e \sim e'$  if and only if  $e$  and  $e'$  are quadrangularly connected in  $G$ . Then  $\sim$  is transitive and an equivalence relation. Thus  $G$  is quadrangularly connected if and only if  $E(G)$  has only one equivalence class.

**Proposition 2.5.** *Every connected  $N_2$ -locally connected graph  $G$  is quadrangularly connected.*

**Proof.** Let  $e_1$  and  $e_2$  be any two edges in  $E(G)$ . If  $e_1$  and  $e_2$  are both in the neighborhood of the second type of  $v$  in  $G$  or  $v$  is a common vertex of  $e_1$  and  $e_2$  (i.e.,  $e_1, e_2 \in E(G[N_2(v) \cup \{v\}])$ ), then it is easy to check that  $e_1 \sim e_2$  since every edge of  $E(G[N_2(v) \cup \{v\}])$  is contained in a 3-cycle or a 4-cycle in  $G[N_2(v) \cup \{v\}]$  from the connectedness of  $N_2(v)$ . Hence  $e_1$  and  $e_2$  belong to two distinct second type neighborhoods of two distinct vertices  $v_1$  and  $v_2$  such that  $e_1 \in E(G[N_2(v_1) \cup \{v_1\}])$  and  $e_2 \in E(G[N_2(v_2) \cup \{v_2\}])$ . Since  $G$  is connected, there is a path  $P = x_0x_1x_2 \dots x_k$  connecting  $N_2(v_1)$  and  $N_2(v_2)$  such that  $x_0, x_1 \in N_2(v_1) \cup \{v_1\}$  and  $x_{k-1}, x_k \in N_2(v_2) \cup \{v_2\}$  but  $x_i$  ( $i = 2, \dots, k - 2$ ) are not in  $N_2(v_1) \cup N_2(v_2) \cup \{v_1, v_2\}$ . For  $i = 1, 2, \dots, k - 1$ , the edges  $x_{i-1}x_i$  and  $x_i x_{i+1}$  are in the same neighborhood of the second type of  $x_i$  in  $G$ , so the edges  $x_{i-1}x_i$  and  $x_i x_{i+1}$  are quadrangularly connected. Note that  $e_1 \sim x_0x_1$  and  $e_2 \sim x_{k-1}x_k$ . Thus, we know from transitivity of  $\sim$  that  $e_1 \sim e_2$ . Therefore, we have completed the proof of Proposition 2.5.  $\square$

In this paper, we show the following result which generalizes Theorems 2.1–2.3. We postpone its proof to the next section.

**Theorem 2.6.** *Every quadrangularly connected claw-free graph without vertices of degree 1, which does not contain an induced subgraph  $H$  isomorphic to either  $G_1$  or  $G_2$  such that  $N_1(x, G)$  of every vertex  $x$  of degree 4 in  $H$  is disconnected is hamiltonian.*

The following was conjectured by Ryjáček [7] and recently proved affirmatively in [3].

**Theorem 2.7** (Lai et al. [3]). *Every 3-connected  $N_2$ -locally connected claw-free graph is hamiltonian.*

We naturally ask whether we may replace  $N_2$ -locally connectedness by quadrangularly connectedness in Theorem 2.7? and so make the following conjecture which would generalize Theorem 2.7 if it is true. Let  $C_1 = (x_1x_2 \dots x_8x_1)$  and  $C_2 = (y_1y_2 \dots y_8y_1)$  be two 8-cycles, and let  $H$  be a graph whose vertex set is  $V(C_1) \cup V(C_2)$  and whose edge set is  $E(C_1) \cup E(C_2) \cup \{x_iy_i : i = 1, 2, \dots, 8\} \cup \{x_1x_5\}$ . Then  $H$  is triangle-free, and the edge  $x_1x_5$  is not contained in a 4-cycle in  $H$ . The line graph  $L(H)$  of  $H$  is 3-connected, quadrangularly connected and claw-free. This example may be useful for considering the following Conjecture 2.8.

**Conjecture 2.8.** *Every 3-connected quadrangularly connected claw-free graph is hamiltonian.*

**Conjecture 2.9.** *Every graph satisfying the conditions of Theorem 2.6 is vertex pancyclic.*

### 3. Proof of Theorem 2.6

In this section, we will provide the proof of Theorem 2.6. Let  $S_{34}(G)$  denote the graph whose vertex set is the set of all 3-cycles and chord-free 4-cycles, and  $C_1, C_2 \in V(C_{34}(G))$  are adjacent in  $S_{34}(G)$  if  $E(C_1) \cap E(C_2) \neq \emptyset$ . Then, from the definition of quadrangularly connectedness, we have established the following fact.

**Proposition 3.1.** *A graph  $G$  is quadrangularly connected if and only if both of the following hold:*

- (1) *For any edge  $e \in E(G)$ , there is a cycle  $C_e \in V(S_{34}(G))$  such that  $e \in E(C_e)$ , and*
- (2)  *$S_{34}(G)$  is connected.*

In our proof, we need to use the Ryjáček's closure concept in claw-free graphs [8]. Let  $G$  be a graph such that for every  $x \in V(G)$ ,  $G[N(x)]$  is either a clique or an union of two disjoint cliques. We call such a graph a *closed claw-free graph*. Then by Lemma 1 in [8], we know that  $G$  is the line graph of a triangle-free graph  $H$ .

**Theorem 3.2** (Ryjáček [8]). *If  $G$  is a claw-free graph, then there is a closed claw-free graph  $\text{cl}(G)$  (called the closure of  $G$ ) such that*

- (1)  *$G$  is a spanning subgraph of  $\text{cl}(G)$ , and*
- (2) *the length of a longest cycle in both  $G$  and  $\text{cl}(G)$  is the same.*

The following fact is easy to prove, and will be used in our proof.

**Proposition 3.3.** *If a claw-free graph  $G$  is quadrangularly connected, then so is  $\text{cl}(G)$ .*

**Proof.** Let  $G$  be a quadrangularly connected claw-free graph and  $v \in V(G)$ . Let  $z, z' \in N_G(v)$  be two nonadjacent vertices in  $G$ , and let  $e$  denote an edge not in  $G$  which joins  $z$  and  $z'$ . Then, in  $G + e$ ,  $e$  lies in a triangle containing  $vz$ , and so  $e \sim vz$  in  $G + e$ . For any edge  $e' \in E(G)$ , since  $G$  is quadrangularly connected,  $vz \sim e'$  in  $G$ , and so  $vz \sim e'$  in  $G + e$ . As  $\sim$  is transitive,  $e \sim e'$  in  $G + e$ . Therefore, we have proved Proposition 3.3.  $\square$

**Proof of Theorem 2.6.** Let  $G$  be a graph satisfying the conditions of Theorem 2.6. If  $G$  is not hamiltonian, then consider the closure  $\text{cl}(G)$  of  $G$ . By Proposition 3.3,  $\text{cl}(G)$  is also quadrangularly connected. From Theorem 3.2(2), without loss of generality assume that  $\text{cl}(G) = G$ . Then  $G$  is a closed claw-free graph. Let  $C$  be a longest cycle of  $G$ . Then there is a vertex  $v$  such that  $v \notin V(C)$  and  $v$  is adjacent to some vertex  $x$  on  $C$ . Let

$$\Gamma = \{e \in E(G) : e \text{ is incident with exactly one vertex in } V(C)\}.$$

Then  $xv \in \Gamma$ . Note that, for any edge  $e \in \Gamma$  (for example,  $e = xv$ ,  $x \in V(C)$  and  $v \notin V(C)$ ), we have that  $x^+v, x^-v \notin E(G)$  since otherwise  $G$  has a longer cycle than  $C$ . Thus  $x^+x^- \in E(G)$  since  $G[x, x^+, x^-, v] \neq K_{1,3}$  (the first vertex  $x$  in the set  $\{x, x^+, x^-, v\}$  is always the center of a claw in the following proof), and we have the following fact.

**Claim 1.** *For any edge  $e \in \Gamma$ ,  $e$  is contained in a 3-cycle  $C_0 \in S_{34}(G)$ .*

**Proof.** Let  $e_1$  be any edge of  $\Gamma$ , and without loss of generality assume that  $e_1 = xv$ ,  $x \in V(C)$  and  $v \notin V(C)$ . Consider the pair of edges  $vx = e_1$  and  $xx^+ = e_2$ . From the conditions of Theorem 2.6, there is a sequence of 3-cycles or chord-free 4-cycles  $C_0, C_1, \dots, C_k$  such that  $e_1 \in E(C_0)$  and  $e_2 \in E(C_k)$  and  $E(C_i) \cap E(C_{i+1}) \neq \emptyset$  for  $i = 0, 1, \dots, k$ . If  $|V(C_0)| = 3$ , we are done. Thus  $|V(C_0)| = 4$  and let  $C_0 = (xvx'x''x)$ .  $\square$

Since  $x^+v, x^-v \notin E(G)$ ,  $e_1$  and  $e_2$  are not contained in some same 3-cycle. We have that  $e_1$  and  $e_2$  are not contained in the same 4-cycle  $C'$  since otherwise, let  $C' = (vxx^+yv)$ . Then  $y \in V(C)$  since otherwise  $G$  has a longer cycle than  $C$  by replacing  $xx^+$  with  $xvyx^+$ , and so  $y^+v, y^-v \notin E(G)$ . It follows that  $y^+y^- \in E(G)$  since  $G$  is claw-free and  $G[y, y^+, y^-, v] \neq K_{1,3}$ , and then  $G$  has a longer cycle than  $C$  by replacing  $y^-yy^+$  with the edge  $y^-y^+$  and  $xx^+$  with  $xvyx^+$ , a contradiction. Thus,  $e_1$  and  $e_2$  are not contained in some same 4-cycle. In order to prove Claim 1, we first establish the following fact.

**Claim 1.1.**  $E(C_0) \cap E(C) = \emptyset$ .

**Proof.** Otherwise,  $|E(C_0) \cap E(C)| = 1$ . We have that  $x'x'' \in E(C)$  because  $xx^+$  and  $xv$  are not at the same cycle in  $S_{34}(G)$ . Thus we can obtain a longer cycle than  $C$  by replacing  $x^-xx^+$  with  $x^-x^+$  and  $x'x''$  with  $x'vx''$ , a contradiction. So Claim 1.1 is true.  $\square$

**Claim 1.2.**  $x'' \in V(C)$ , and  $x^-x'', x^+x'' \in E(G)$ .

**Proof.** If  $x'' \notin V(C)$ , then  $xx'' \in \Gamma$ , and so  $x''x^+, x''x^- \notin E(G)$ . Note that  $x''v \notin E(G)$ . Thus  $G[x, v, x'', x^+] = K_{1,3}$ , a contradiction. Since  $G[x, x^-(x^+), v, x''] \neq K_{1,3}$  and  $vx'', vx^+, vx^- \notin E(G)$ ,  $x^-x'', x^+x'' \in E(G)$ .  $\square$

**Claim 1.3.**  $x' \in V(C)$  and  $x'^+x'^- \in E(G)$ .

**Proof.** If  $x' \notin V(C)$ , then  $x'x'' \in \Gamma$ , and so  $x''^-x''^+ \in E(G)$  since  $G[x'', x''^+, x''^-, x'] \neq K_{1,3}$ . It follows that we can obtain a longer cycle than  $C$  by replacing  $x''^-x''^+x''^+$  with  $x''^-x''^+$  and  $xx^+$  with  $xvx'x''x^+$ , a contradiction. Thus  $x' \in V(C)$ , and so  $vx' \in \Gamma$ . From  $G[x', x'^+, x'^-, v] \neq K_{1,3}$  and  $x'^+v, x'^-v \notin E(G)$ , we have  $x'^+x'^- \in E(G)$ . Thus Claim 1.3 is true.  $\square$

Now we return the proof of Claim 1. By Claims 1.2 and 1.3, we have  $x', x'' \in V(C)$ . Without loss of generality assume that  $x'' \in C(x, x')$ . By Claim 1.1,  $|C(x'', x')| \geq 1$ . From  $G[x', v, x'^+(x'^-), x''] \neq K_{1,3}$  and  $x'^+v, x'^-v, x''v \notin E(G)$ ,  $x'^+x'', x'^-x'' \in E(G)$ . Similarly,  $x^+x'', x^-x'' \in E(G)$ . It follows that we easily get  $|C(x^+, x'')| \geq 1$  and  $|C(x'', x'^-)| \geq 1$ . Note that  $xx' \notin E(G)$ . From  $G[x'', x''^-, x, x'] \neq K_{1,3}$ ,  $x''^-x \in E(G)$  or  $x''^-x' \in E(G)$ . If  $x''^-x \in E(G)$ , then  $G$  has a new cycle  $C[x^+, x''^-]xvx'C[x'', x'^-]C[x'^+, x^-]x^+$  than  $C$ , a contradiction. Similarly, if  $x''^-x' \in E(G)$ , we easily get a longer cycle than  $C$ , a contradiction. Thus we have completed the proof of Claim 1.  $\square$

By Proposition 3.1(1), there is a sequence of  $C_1, C_2, \dots, C_m$  in  $S_{34}(G)$  such that  $E(C)$  is contained in  $E(C_1) \cup E(C_2) \cup \dots \cup E(C_m)$  and  $m$  is minimal with the property. Then, by the minimality of  $m$ , we have the following fact.

**Claim 2.**  $E(C_i) \cap E(C) \neq \emptyset$  for  $i = 1, 2, \dots, m$ .

For any edge  $e \in \Gamma$ , by Claim 1, there is a 3-cycle  $C_0$  such that  $e \in E(C_0)$ . By Proposition 3.1, there is a shortest path  $P = Q_0Q_1Q_2 \dots Q_r$  in  $S_{34}(G)$  from  $C_0$  to the vertex set  $\{C_1, C_2, \dots, C_m\}$ , where  $Q_0 = C_0$ . From the above, we know that the length  $r$  of  $P$  is at least one. Choose an edge (say  $e$ ) in  $\Gamma$  such that  $r$  is as small as possible. without loss of generality assume that  $e = xv$  and  $Q_r = C_1$ . Then we have the following fact.

**Claim 3.**  $|V(C_0) \cap V(C)| \leq 2$ , and  $E(Q_i) \cap E(C) = \emptyset$  for  $i = 0, 1, \dots, r - 1$ .

**Claim 4.**  $r = 1$  and  $|V(C_0) \cap V(C)| = 2$ .

**Proof.** If  $r > 1$ , then, since  $|V(C_0) \cap V(C)| \leq 2$ , there is a largest integer  $t$  ( $0 \leq t \leq r - 1$ ) such that  $|V(Q_t) \cap V(C)| = 2$ , and by Claim 1,  $|V(Q_t)| = 3$ . If  $t = 0$ , then we are done. Thus  $t \geq 1$ . Let  $Q_t = (v'u'w'v')$  and  $v' \notin V(C)$  but  $u' \in V(C)$ . We replace the edge  $e$  by  $v'u'$ , we get a shorter path  $P' = Q_t \dots Q_r$  than  $P$ , which contradicts the choice of  $e$ . Thus  $r = 1$ . From the above, we also have  $|V(C_0) \cap V(C)| = 2$ .  $\square$

Let  $C_0 = (xvwx)$ . Then  $w, x \in V(C)$ , and so  $vw \in \Gamma$  and  $w^+w^- \in E(G)$  but  $w^-v, w^+v \notin E(G)$ . Note that  $Q_1 = C_1$  and  $C$  is divided into two segments  $C' = (xx^+ \dots w^-w)$  and  $C'' = (ww^+ \dots x^-x)$ . We further have the following fact.

**Claim 5.**  $|V(C_1)| = 4$ .

**Proof.** Otherwise,  $|V(C_1)| = 3$ . Since  $E(C) \cap E(C_1) \neq \emptyset$  and  $E(C_0) \cap E(C_1) \neq \emptyset$ , we must have that one of  $\{ww^-, ww^+, xx^+, xx^-\}$  belongs to  $E(C) \cap E(C_1)$ . For example,  $ww^- \in E(C) \cap E(C_1)$ . Since  $vw^- \notin E(G)$ , we have  $C_1 = (xww^-x)$ . Note that  $w^+w^-, x^+x^- \in E(G)$ . Thus replacing  $x^-xx^+$  by  $x^-x^+$  and  $w^-w$  by  $w^-xvw$ , we obtain a longer than  $C$ , a contradiction. Thus Claim 5 is true.  $\square$

Let  $e_1 = yz \in E(C) \cap E(C_1)$ . Then  $e_1 \in \{xx^+, x^-x, w^-w, ww^+\}$  since otherwise, for example,  $yz$  is on the segment  $C(x^+, w^-)$ . Since  $E(C_0) \cap E(C_1) \neq \emptyset$ , there is an edge  $e'' (\neq e_1)$  such that  $e'' \in E(C_0) \cap E(C_1)$  and

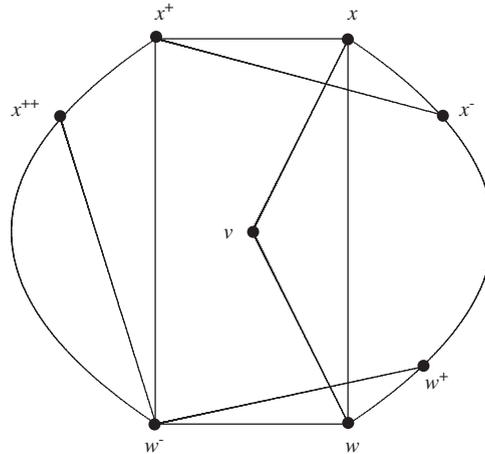


Fig. 2.

$e'' \in \{xv, vw, xw\}$ . If  $e'' = xv$ , then  $vz \in \Gamma$  and without loss of generality assume that  $C_1 = (xvzyx)$ . Replacing  $yz$  by  $yxvz$  and  $x^-xx^+$  by  $x^-x^+$ , we obtain a longer cycle than  $C$ , a contradiction. Similarly, if  $e'' = vw$ , then we can get a contradiction. If  $e'' = xw$ , then replacing  $yz$  with  $yxvwz$ ,  $x^-xx^+$  with  $x^-x^+$  and  $w^-ww^+$  with  $w^-w^+$ , we get a longer cycle than  $C$ . This contradiction shows  $e_1 \in \{xx^+, x^-x, w^-w, ww^+\}$  (say  $e_1 = w^-w$ ).

Note that  $vw^- \notin E(G)$ . If  $vw \in E(C_0) \cap E(C_1)$ , let  $C_1 = (wvzw^-w)$ . Then  $z \in V(C)$  and so  $vz \in \Gamma$  and  $z^+z^- \in E(G)$ . Without loss of generality assume that  $z \in C(x^+, w^-)$ , we obtain a longer cycle than  $C$  by replacing  $z^-zz^+$  with  $z^-z^+$  and  $w^-w$  with  $w^-zvw$ . This contradiction shows  $vw \notin E(C_0) \cap E(C_1)$ . Similarly,  $vx \notin E(C_0) \cap E(C_1)$ . Thus  $xw \in E(C_0) \cap E(C_1)$ .  $\square$

Let  $C_1 = (xww^-zx)$ . Then we have the following fact.

**Claim 6.** Without loss of generality, we may assume that  $z = x^+$ , and  $C_1 = (xww^-x^+x)$ .

**Proof.** Otherwise, without loss of generality assume that  $z \in C(x^+, w^-)$ . Then  $|C(x^+, z)| \geq 1$ . We have  $z^-z^+ \notin E(G)$  since otherwise we can obtain a longer cycle than  $C$  by replacing  $z^-zz^+$  with  $z^-z^+$ ,  $x^-xx^+$  with  $x^-x^+$ , and  $w^-w$  with  $w^-zxvw$ . Obviously,  $z^-w^- \notin E(G)$  since otherwise  $G$  has a longer cycle  $C[x^+, z^-]C^-[w^-, z]xvC[w, x^-]x^+$  than  $C$ . From  $G[z, x, z^-, w^-] \neq K_{1,3}$ , we have  $xz^- \in E(G)$ . Replacing  $x^-xx^+$  with  $x^-x^+$  and  $z^-z$  with  $z^-xz$ , we obtain a new cycle  $C'$  of the same length as  $C$  but  $xz$  is an edge on  $C'$ . Thus without loss of generality, we may assume that  $z = x^+$ , and so  $C_1 = (xww^-x^+x)$ . Therefore, Claim 6 is true.  $\square$

We next complete the proof of Theorem 2.6. By Claims 5 and 6, we have  $x^+x^-, w^-w^+ \in E(G)$ . It is easy to see that  $|C(x^+, w^-)| \geq 1$  and  $|C(w^+, x^-)| \geq 1$ . Let  $x^{++} = (x^+)^+$ . Obviously,  $x^-x^{++} \notin E(G)$  since otherwise we can obtain a longer cycle by replacing  $x^-xx^+x^{++}$  with  $x^-x^{++}$  and  $w^-w$  with  $w^-x^+xvw$ . We have  $x^-w^- \notin E(G)$  since otherwise  $G$  has a longer cycle  $C[x, w^-]C^-[x^-, w]vx$  than  $C$ . Similarly,  $x^+w^+ \notin E(G)$ . From  $G[x^+, x^-, x^{++}, w^-] \neq K_{1,3}$ , we have  $x^{++}w^- \in E(G)$  (see Fig. 2). Let  $H = G[V(C_1) \cup \{x^{++}, w^+, v, x^-\}]$  (see Fig. 2). Since  $G$  is a closed claw-free graph, it is easy to check that  $G[N(x)]$ ,  $G[N(x^+)]$ ,  $G[N(w^-)]$  and  $G[N(w)]$  are not connected. We have that  $vx^{++} \notin E(G)$  since otherwise  $G$  has a longer cycle  $C[x^{++}, x^-]x^+vx^{++}$  than  $C$ .

If  $x^-w^- \in E(G)$ , then  $G$  has a longer cycle  $C[w, x^-]w^-C^-[w^-, x]vw$  than  $C$ . This contradiction shows that  $x^-w^- \notin E(G)$ . Similarly, we can prove that

$$x^-w^-, w^+x^{++}, x^-x^{++}, vx^-, w^+v, w^+x^- \notin E(G).$$

Note that it is possible that  $x^-w^+ \in E(G)$ . Thus  $H$  is isomorphic to  $G_1$  or  $G_2$  in Fig. 1, a contradiction and the proof of Theorem 2.6 has been completed.  $\square$

## References

- [1] J.A. Bondy, U.S.R. Murty, Graph Theory with Applications, Macmillan, Elsevier, New York, London, 1976.
- [2] Y. Shao, Claw-Free Graphs and Line Graphs, Ph.D. Dissertation, West Virginia University 2005.
- [3] H.-J. Lai, Y. Shao, M. Zhan, Hamiltonian  $N_2$ -locally connected claw-free graphs, J. Graph Theory 48 (2005) 142–146.
- [4] M. Li, On pancyclic claw-free graphs, Ars Combinatoria 50 (1998) 279–291.
- [5] M. Li, Hamiltonian cycles in  $N^2$ -locally connected claw-free graphs, Ars Combinatoria 62 (2002) 281–288.
- [6] D.J. Oberly, D.P. Sumner, Every connected, locally connected nontrivial graph with no induced claw is hamiltonian, J. Graph Theory 3 (1979) 351–356.
- [7] Z. Ryjáček, Hamiltonian circuits in  $N_2$ -locally connected  $K_{1,3}$ -free graphs, J. Graph Theory 14 (1990) 321–331.
- [8] Z. Ryjáček, On a closure concept in claw-free graphs, J. Combin. Theory Ser. B 70 (1997) 217–224.